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EXAMINER				
LOUTE, OSCAR A				
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Please find below and/or attached an Office communication concerning this application or proceeding.

The time period for reply, if any, is set in the attached communication.

Office Action Summary

Application No.

10/824,729

Applicant(s)

BAUGHER, MARK

Examiner

OSCAR A. LOUIE

Art Unit

2136

-- The MAILING DATE of this communication appears on the cover sheet with the correspondence address --
Period for Reply

A SHORTENED STATUTORY PERIOD FOR REPLY IS SET TO EXPIRE 3 MONTH(S) OR THIRTY (30) DAYS, WHICHEVER IS LONGER, FROM THE MAILING DATE OF THIS COMMUNICATION.

- Extensions of time may be available under the provisions of 37 CFR 1.136(a). In no event, however, may a reply be timely filed after SIX (6) MONTHS from the mailing date of this communication.
- If NO period for reply is specified above, the maximum statutory period will apply and will expire SIX (6) MONTHS from the mailing date of this communication.
- Failure to reply within the set or extended period for reply will, by statute, cause the application to become ABANDONED (35 U.S.C. § 133). Any reply received by the Office later than three months after the mailing date of this communication, even if timely filed, may reduce any earned patent term adjustment. See 37 CFR 1.704(b).

Status

- 1) ☒ Responsive to communication(s) filed on 23 May 2008.
- 2a) ☐ This action is **FINAL**. 2b) ☒ This action is non-final.
- 3) ☐ Since this application is in condition for allowance except for formal matters, prosecution as to the merits is closed in accordance with the practice under *Ex parte Quayle*, 1935 C.D. 11, 453 O.G. 213.

Disposition of Claims

- 4) ☒ Claim(s) 1-31 is/are pending in the application.
- 4a) Of the above claim(s) _____ is/are withdrawn from consideration.
- 5) ☐ Claim(s) _____ is/are allowed.
- 6) ☒ Claim(s) 1-31 is/are rejected.
- 7) ☐ Claim(s) _____ is/are objected to.
- 8) ☐ Claim(s) _____ are subject to restriction and/or election requirement.

Application Papers

- 9) ☐ The specification is objected to by the Examiner.
- 10) ☐ The drawing(s) filed on _____ is/are: a) ☐ accepted or b) ☐ objected to by the Examiner.
Applicant may not request that any objection to the drawing(s) be held in abeyance. See 37 CFR 1.85(a).
Replacement drawing sheet(s) including the correction is required if the drawing(s) is objected to. See 37 CFR 1.121(d).
- 11) ☐ The oath or declaration is objected to by the Examiner. Note the attached Office Action or form PTO-152.

Priority under 35 U.S.C. § 119

- 12) ☐ Acknowledgment is made of a claim for foreign priority under 35 U.S.C. § 119(a)-(d) or (f).
- a) ☐ All b) ☐ Some * c) ☐ None of:
1. ☐ Certified copies of the priority documents have been received.
 2. ☐ Certified copies of the priority documents have been received in Application No. _____.
 3. ☐ Copies of the certified copies of the priority documents have been received in this National Stage application from the International Bureau (PCT Rule 17.2(a)).

* See the attached detailed Office action for a list of the certified copies not received.

Attachment(s)

- 1) ☒ Notice of References Cited (PTO-892)
- 2) ☐ Notice of Draftsperson's Patent Drawing Review (PTO-948)
- 3) ☐ Information Disclosure Statement(s) (PTO/SF/ICE)
Paper No(s)/Mail Date _____
- 4) ☐ Interview Summary (PTO-413)
Paper No(s)/Mail Date _____
- 5) ☐ Notice of Informal Patent Application
- 6) ☐ Other: _____

DETAILED ACTION

This first non-final action is in response to the Request for Continued Examination filing of 05/23/2008. Claims 1-31 are pending and have been considered as follows.

Examiner Note

In light of the applicant's amendments, the examiner hereby withdraws his previous Specification Objection and his previous 35 U.S.C. 101 rejections with respect to Claims 18 & 19.

Claim Objections

1. Claim 18 is objected to because of the following informalities:
 - Claim 18 line 5 recites the term "when" which should be omitted;Appropriate correction is required.

Claim Rejections - 35 USC § 101

2. 35 U.S.C. 101 reads as follows:

Whoever invents or discovers any new and useful process, machine, manufacture, or composition of matter, or any new and useful improvement thereof, may obtain a patent therefor, subject to the conditions and requirements of this title.

Claim 17 is rejected under 35 U.S.C. 101 because the claimed invention is directed to non-statutory subject matter.

- Claim 17 recites “an apparatus” comprising “means for” where the “means for” appear to be nothing more than computer program modules, thereby invoking 35 U.S.C. 101 as non-statutory subject matter;

Claim Rejections - 35 USC § 102

3. The following is a quotation of the appropriate paragraphs of 35 U.S.C. 102 that form the basis for the rejections under this section made in this Office action:

A person shall be entitled to a patent unless –

(e) the invention was described in (1) an application for patent, published under section 122(b), by another filed in the United States before the invention by the applicant for patent or (2) a patent granted on an application for patent by another filed in the United States before the invention by the applicant for patent, except that an international application filed under the treaty defined in section 351(a) shall have the effects for purposes of this subsection of an application filed in the United States only if the international application designated the United States and was published under Article 21(2) of such treaty in the English language.

4. Claims 17-19 are rejected under 35 U.S.C. 102(e) as being anticipated by Schuba et al. (US-6944663-B2).

Claims 17-19:

Schuba et al. disclose an apparatus for preventing an attack on a network and a computer-readable volatile or non-volatile medium storing one or more sequences of instructions comprising,

- “a processor” (i.e. “a computer system based on a microprocessor, a mainframe computer, a digital signal processor, a portable computing device, a personal organizer, a device controller, and a computational engine within an appliance”) [column 3 lines 35-38];

- “a computer-readable volatile or non-volatile medium having- stored thereon one or more stored sequences of instructions that are accessible to the processor” (i.e. “The data structures and code described in this detailed description are typically stored on a computer readable storage medium, which may be any device or medium that can store code and/or data for use by a computer system. This includes, but is not limited to, magnetic and optical storage devices such as disk drives, magnetic tape, CDs (compact discs) and DVDs (digital versatile discs or digital video discs)”) [column 3 lines 10-12];
- “receiving a request to access a resource from a user” (i.e. “the system receives a request for service from a client 106 (step 202)”) [column 3 lines 52-53];
“wherein the request includes an accumulated work value” (i.e. “the system generates a random number, y , and a transaction identifier, $id.sub.1$ (step 204). The system also selects a value for the parameter, n , which specifies the amount of computational work involved in computing the preimage x , such that $h(x)=y$ (step 206)”) [column 3 lines 53-58];
- “determining whether the accumulated work value exceeds a required work threshold value” (i.e. “If $id.sub.1 = id.sub.2$ at step 218, the system computes $h(x)$ (step 220). Next, the system compares y and $h(x)$ (step 222). If $y=h(x)$, the client successfully solved the client puzzle, and the system performs the requested service for the client (step 224)”) [column 4 lines 35-39];

- “if not, selectively requiring the user to perform a quantity of work as a condition for accessing the resource” (i.e. “FIG. 2 is a flow chart illustrating the process of using a client puzzle in accordance with an embodiment of the present invention”) [column 3 lines 50-52];
- “providing the user with access to the resource” (i.e. “Next, the system compares y and $h(x)$ (step 222). If $y=h(x)$, the client successfully solved the client puzzle, and the system performs the requested service for the client (step 224)”) [column 4 lines 36-39];
- “determining an amount of accumulated work output value to provide to the user based on a volume of data communicated between the resource and the user” (i.e. “the system generates a random number, y , and a transaction identifier, $id.sub.1$ (step 204). The
- system also selects a value for the parameter, n , which specifies the amount of computational work involved in computing the preimage x , such that $h(x)=y$ (step 206)”) [column 3 lines 53-58];
- “providing the accumulated work output value to the user” (i.e. “The system also selects a value for the parameter, n , which specifies the amount of computational work involved in computing the preimage x , such that $h(x)=y$ (step 206)”) [column 3 lines 55-58].

Claim Rejections - 35 USC § 103

5. The following is a quotation of 35 U.S.C. 103(a) which forms the basis for all obviousness rejections set forth in this Office action:

(a) A patent may not be obtained though the invention is not identically disclosed or described as set forth in section 102 of this title, if the differences between the subject matter sought to be patented and the prior art are such that the subject matter as a whole would have been obvious at the time the invention was made to a person having ordinary skill in the art to which said subject matter pertains. Patentability shall not be negated by the manner in which the invention was made.

6. Claims 20, 24, & 28 are rejected under 35 U.S.C. 103(a) as being unpatentable over Schuba et al. (US-6944663-B2).

Claims 20, 24, & 28:

Schuba et al. disclose an apparatus for preventing an attack on a network and a computer-readable volatile or non-volatile medium storing one or more sequences of instructions, as in Claims 17-19 above, further comprising,

- “determining whether a mathematical relationship of the current user identity value and the prior user identity value indicates that the user has possession of a resource secret” (i.e. “If $id.sub.1 = id.sub.2$ at step 218, the system computes $h(x)$ (step 220). Next, the system compares y and $h(x)$ (step 222). If $y=h(x)$, the client successfully solved the client puzzle, and the system performs the requested service for the client (step 224)”) [column 4 lines 35-39].

but they do not explicitly disclose,

- “wherein the request includes a prior user identity value and a current user identity value,” although Schuba et al. do suggest two separate identifiers, as recited below;

however, they do disclose,

- “For example, if the parameters associated with the client (id.sub.1, n, y) are stored in a database that is indexed by id.sub.1, a subsequent lookup using id.sub.2 will return (id.sub.1, n, y) only if id.sub.1 = id.sub.2. Alternatively, if the lookup is based on client identifiers, an explicit comparison of id.sub.1 and id.sub.2 needs to be performed”
[column 4 lines 29-34];

Therefore, it would have been obvious for one of ordinary skill in the art at the time of the applicant’s invention to include, “wherein the request includes a prior user identity value and a current user identity value,” in the invention as disclosed by Schuba et al. since it would be expected that a client/user may attempt to connect more than just once and accommodations need to be made to handle the scenarios where the client is legitimate and non-legitimate as is suggested by Schuba et al.

7. Claims 1, 2, 11, 15, & 16 are rejected under 35 U.S.C. 103(a) as being unpatentable over Schuba et al. (US-6944663-B2) in view of French et al. (US-6321339-B1).

Claim 1:

Schuba et al. disclose a method of preventing an attack on a network comprising,

- “receiving a request to access a resource from a user” (i.e. “the system receives a request for service from a client 106 (step 202)”) [column 3 lines 52-53];

- “wherein the request includes an accumulated work value” (i.e. “the system generates a random number, y , and a transaction identifier, $id.sub.1$ (step 204). The system also selects a value for the parameter, n , which specifies the amount of computational work involved in computing the preimage x , such that $h(x)=y$ (step 206)”) [column 3 lines 53-58];
- “determining whether the accumulated work value exceeds a required work threshold value” (i.e. “If $id.sub.1 = id.sub.2$ at step 218, the system computes $h(x)$ (step 220). Next, the system compares y and $h(x)$ (step 222). If $y=h(x)$, the client successfully solved the client puzzle, and the system performs the requested service for the client (step 224)”) [column 4 lines 35-39];
- “if not, requiring the user to perform a quantity of work as a condition for accessing the resource” (i.e. “FIG. 2 is a flow chart illustrating the process of using a client puzzle in accordance with an embodiment of the present invention”) [column 3 lines 50-52];
- “providing the user with access to the resource” (i.e. “Next, the system compares y and $h(x)$ (step 222). If $y=h(x)$, the client successfully solved the client puzzle, and the system performs the requested service for the client (step 224)”) [column 4 lines 36-39];
- “determining an amount of accumulated work output value to provide to the user based on a volume of data communicated between the resource and the user” (i.e. “the system generates a random number, y , and a transaction identifier, $id.sub.1$ (step 204). The system also selects a value for the parameter, n , which specifies the amount of computational work involved in computing the preimage x , such that $h(x)=y$ (step 206)”) [column 3 lines 53-58];

- “providing the accumulated work output value to the user” (i.e. “The system also selects a value for the parameter, n , which specifies the amount of computational work involved in computing the preimage x , such that $h(x)=y$ (step 206)”) [column 3 lines 55-58];

but, they do not disclose,

- “wherein the accumulated work value represents a total amount of work previously performed by the user and accumulated across multiple prior requests,” although French et al. do suggest combined authentication scores, as recited below;
- “wherein the accumulated work output value represents a second amount of work performed by the user,” although French et al. do suggest combined authentication scores, as recited below;

however, French et al. do disclose,

- “The transaction record 112 (illustrated in FIGS. 13-16) initialized in step 22 is used throughout the authentication process 10 to keep track of user input and authentication results. After the appropriate queries have been processed and all results stored in the transaction record 112, the transaction record 112 is used to determine an authentication score with respect to the request. Step 56 calculates the total authentication score, and optionally, a score for each data source, data field, etc. The results are categorized as a big hit (B), a regular hit (R), a possible hit (P), or no hit (N) depending on results. Those results may then be combined with the results of second level authentication process 40 to determine an overall authenticity certainty score, as illustrated in FIGS. 23-28 and discussed below” [column 14 lines 1-14];

Therefore, it would have been obvious for one of ordinary skill in the art at the time of the applicant's invention to include, "wherein the accumulated work value represents a total amount of work previously performed by the user and accumulated across multiple prior requests" and "wherein the accumulated work output value represents a second amount of work performed by the user," in the invention as disclosed by Schuba et al. for the purposes of providing categories of access based on values associated with authentication outcomes.

Claim 2:

Schuba et al. and French et al. disclose a method of preventing an attack on a network, as in Claim 1 above, their combination further disclosing,

- "determining whether a mathematical relationship of the current user identity value and the prior user identity value indicates that the user has possession of a resource secret" (i.e. "If $\text{id.sub.1} = \text{id.sub.2}$ at step 218, the system computes $h(x)$ (step 220). Next, the system compares y and $h(x)$ (step 222). If $y=h(x)$, the client successfully solved the client puzzle, and the system performs the requested service for the client (step 224)") [column 4 lines 35-39].

but they do not explicitly disclose,

- "wherein the request includes a prior user identity value and a current user identity value," although Schuba et al. do suggest separate identifiers, as recited below;

however, they do disclose,

- “For example, if the parameters associated with the client (id.sub.1, n, y) are stored in a database that is indexed by id.sub.1, a subsequent lookup using id.sub.2 will return (id.sub.1, n, y) only if id.sub.1 = id.sub.2. Alternatively, if the lookup is based on client identifiers, an explicit comparison of id.sub.1 and id.sub.2 needs to be performed”
[column 4 lines 29-34];

Therefore, it would have been obvious for one of ordinary skill in the art at the time of the applicant's invention to include, “wherein the request includes a prior user identity value and a current user identity value,” in the invention as disclosed by Schuba et al. and French et al. since it would be expected that a client/user may attempt to connect more than just once and accommodations need to be made to handle the scenarios where the client is legitimate and non-legitimate as is suggested by Schuba et al.

Claim 11:

Schuba et al. and French et al. disclose a method of preventing an attack on a network, as in Claim 1 above, their combination further disclosing,

- “receiving the accumulated proof of work value” (i.e. “If id.sub.1 = id.sub.2 at step 218, the system computes $h(x)$ (step 220). Next, the system compares y and $h(x)$ (step 222). If $y=h(x)$, the client successfully solved the client puzzle, and the system performs the requested service for the client (step 224)”) [column 4 lines 35-39].

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Claim 15:

Schuba et al. and French et al. disclose a method of preventing an attack on a network, as in

Claim 1 above, their combination further disclosing,

- “wherein requiring the user to perform a quantity of work as a condition for accessing the resource comprises requiring the user to hash a message until a specified number of bits are zero” (i.e. “Next, the system stores (id.sub.1, n, y) at server 102 (step 208) and sends (id.sub.1, n, y) to client 106 (step 210). The system then allows client 106 to compute the preimage x, such that $h(x)=y$ (step 212). In one embodiment of the present invention, h is a hash function, such as SHA1 or MD5, so that computing the preimage x given y requires significantly more time than computing the hash function $h(x)$ given x”) [column 3 lines 59-64].

Claim 16:

Schuba et al. disclose a method of preventing an attack on a network comprising,

- “receiving a request to access a resource from a user” (i.e. “the system receives a request for service from a client 106 (step 202).”) [column 3 lines 52-53];
- “determining whether the accumulated work value exceeds a required work threshold value” (i.e. “If id.sub.1 =id.sub.2 at step 218, the system computes $h(x)$ (step 220). Next, the system compares y and $h(x)$ (step 222). If $y=h(x)$, the client successfully solved the client puzzle, and the system performs the requested service for the client (step 224)”) [column 4 lines 35-39];

- “providing the user with access to the resource only when the accumulated work value exceeds a required work threshold value” (i.e. “If $y=h(x)$, the client successfully solved the client puzzle, and the system performs the requested service for the client (step 224)”) [column 4 lines 36-39];

but they do not disclose,

- “wherein the request includes an accumulated work value that represents work that the resource has previously required the user to perform in order to obtain previous access to the resource,” although French et al. do suggest combined authentication scores, as recited below;

however, French et al. do disclose,

- “The transaction record 112 (illustrated in FIGS. 13-16) initialized in step 22 is used throughout the authentication process 10 to keep track of user input and authentication results. After the appropriate queries have been processed and all results stored in the transaction record 112, the transaction record 112 is used to determine an authentication score with respect to the request. Step 56 calculates the total authentication score, and optionally, a score for each data source, data field, etc. The results are categorized as a big hit (B), a regular hit (R), a possible hit (P), or no hit (N) depending on results. Those results may then be combined with the results of second level authentication process 40 to determine an overall authenticity certainty score, as illustrated in FIGS. 23-28 and discussed below” [column 14 lines 1-14];

Therefore, it would have been obvious for one of ordinary skill in the art at the time of the applicant's invention to include, "wherein the request includes an accumulated work value that represents work that the resource has previously required the user to perform in order to obtain previous access to the resource," in the invention as disclosed by Schuba et al. for the purposes of providing categories of access based on values associated with authentication outcomes.

8. Claims 3-10 & 12-14 are rejected under 35 U.S.C. 103(a) as being unpatentable over Schuba et al. (US-6944663-B2) in view of French et al. (US-6321339-B1) and in further view of Juels et al. (US-7197639-B1).

Claim 3:

Schuba et al. and French et al. disclose a method of preventing an attack on a network, as in Claim 1 above, further comprising,

- "wherein $H(i+1,x)$ is computed by the user as a hash chain from a non-shared user secret (x)" (i.e. "Next, the system stores (id.sub.1, n, y) at server 102 (step 208) and sends (id.sub.1, n, y) to client 106 (step 210). The system then allows client 106 to compute the preimage x , such that $h(x)=y$ (step 212). In one embodiment of the present invention, h is a hash function, such as SHA1 or MD5, so that computing the preimage x given y requires significantly more time than computing the hash function $h(x)$ given x ") [column 3 lines 59-64];
- "wherein $H(n,x)=h(H(n-1,x))$ " (i.e. "If id.sub.1=id.sub.2 at step 218, the system computes $h(x)$ (step 220). Next, the system compares y and $h(x)$ (step 222). If $y=h(x)$, the client successfully solved the client puzzle, and the system performs the requested service for the client (step 224)") [column 4 lines 35-39];

- “wherein $n > 0$ and $H(0,x) = x$ ” (i.e. “The parameter n is used to adjust the amount of work required to compute the preimage x . For example, the parameter n can be used as a parameter to the hash function h , which indicates both the size of the hash value generated by the hash function h , as well as the number of bits of x that are used in computing $h(x)$ ”) [column 4 lines 3-8];
- “wherein function h is a one-way function that is difficult to invert” (i.e. “ h is a hash function, such as SHA1 or MD5, so that computing the preimage x given y requires significantly more time than computing the hash function $h(x)$ given x ”) [column 3 lines 63-64];
- “receiving a current user identity value $H(i,x)$ ” (i.e. “Next, the system receives (id.sub.2, x) from the client (step 214), wherein id.sub.2 is an identifier returned by the client and x is the preimage of y computed by the client”) [column 4 lines 20-22];
- “verifying that the keyless user identity value properly identifies the user only upon determining that $h(H(i,x)) = H(i+1,x)$ ” (i.e. “If id.sub.1 = id.sub.2 at step 218, the system computes $h(x)$ (step 220). Next, the system compares y and $h(x)$ (step 222). If $y=h(x)$, the client successfully solved the client puzzle, and the system performs the requested service for the client (step 224)”) [column 4 lines 35-39];

but they do not disclose,

- “receiving a prior keyless user identity value $H(i+1,x)$ in the request comprising a one-time password,” although Juels et al. do suggest a secret password, as recited below;

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however, Juels et al. do disclose,

- “For example, after TCP-IP is established, the next higher protocol layer can demand a secret password or other form of authentication before proceeding with the execution of the server application” [column 13 lines 23-25];

Therefore, it would have been obvious for one of ordinary skill in the art at the time of the applicant’s invention to include, “receiving a prior keyless user identity value $H(i+1,x)$ in the request comprising a one-time password,” in the invention as disclosed by Schuba et al. and French et al. since “an adversary cannot pass through this security barrier. If this were not true, then the adversary would not be limited to disabling the server 120 through session-establishing resource depletion” [column 13 lines 27-30].

Claim 4:

Schuba et al., French et al., and Juels et al. disclose a method of preventing an attack on a network, as in Claim 3 above, their combination further disclosing,

- “wherein h comprises a SHA-1 hash algorithm” (i.e. “h is a hash function, such as SHA1 or MD5, so that computing the preimage x given y requires significantly more time than computing the hash function $h(x)$ given x”) [column 3 lines 63-64].

Claim 5:

Schuba et al., French et al., and Juels et al. disclose a method of preventing an attack on a network, as in Claim 3 above, their combination further disclosing,

- “wherein n is between 10^4 and 10^6 ” (i.e. “The parameter n is used to adjust the amount of work required to compute the preimage x. For example, the parameter n can be used as

a parameter to the hash function h , which indicates both the size of the hash value generated by the hash function h , as well as the number of bits of x that are used in computing $h(x)$ ” [column 4 lines 3-8].

Claim 6:

Schuba et al. and French et al. disclose a method of preventing an attack on a network, as in Claim 3 above, their combination further disclosing, disclose a method of preventing an attack on a network, as in Claim 1 above, but their combination do not disclose,

- “determining the required work threshold value based on a then-current capacity of the resource,” although Juels et al. do suggest computational capacity used to determine computational size, as recited below;

however, Juels et al. do disclose,

- “the rate of connection buffer allocation and the likely computational capacity of one or more attacking clients 110 can be used to select the computational size of a particular tasks when operating in a defensive mode” [column 7 lines 29-33];

Therefore, it would have been obvious for one of ordinary skill in the art at the time of the applicant’s invention to include, “determining the required work threshold value based on a then-current capacity of the resource,” in the invention as disclosed by Schuba et al. and French et al. for the purposes of assessing the likelihood of attack.

Claim 7:

Schuba et al. and French et al. disclose a method of preventing an attack on a network, as in Claim 1 above, but their combination do not disclose,

- “determining the required work threshold value based on a then-current capacity of the resource,” although Juels et al. do suggest computational capacity determining computational size, as recited below;
- “requiring a first user who has an accumulated work value that is greater than the required work threshold value to perform a first amount of work as a condition for accessing the resource,” although Juels et al. do suggest adjusting puzzle size/complexity, as recited below;
- “requiring a second user who has an accumulated work value that is less than or equal to the required work threshold value to perform a second amount of work as a condition for accessing the resource,” although Juels et al. do suggest adjusting puzzle size/complexity, as recited below;
- “wherein the second amount of work is greater than the first amount of work,” although Juels et al. do suggest adjusting puzzle size/complexity, as recited below;

however, Juels et al. do disclose,

- “the rate of connection buffer allocation and the likely computational capacity of one or more attacking clients 110 can be used to select the computational size of a particular tasks when operating in a defensive mode” [column 7 lines 29-33];

- “The client puzzle protocol also allows for graceful degradation in service when an attack is mounted. The size of the puzzles can be increased as the progress of an attack advances closer to disabling the server. This enables the protocol to flex according to the scale of the attack” [column 9 lines 10-14];

Therefore, it would have been obvious for one of ordinary skill in the art at the time of the applicant’s invention to include, “determining the required work threshold value based on a then-current capacity of the resource” and “requiring a first user who has an accumulated work value that is greater than the required work threshold value to perform a first amount of work as a condition for accessing the resource” and “requiring a second user who has an accumulated work value that is less than or equal to the required work threshold value to perform a second amount of work as a condition for accessing the resource” and “wherein the second amount of work is greater than the first amount of work,” in the invention as disclosed by Schuba et al. and French et al. since the client puzzle protocol is used for controlling the rate of connection buffer allocation and the likely computational capacity in order to provide graceful degradation in service when an attack is mounted (i.e. denial of service attack).

Claim 8:

Schuba et al. and French et al. disclose a method of preventing an attack on a network, as in

Claim 1 above, but their combination do not disclose,

- “wherein the step of determining an amount of accumulated work output value is performed for a specified user only during a specified time period in which accumulating work is allowed for that specified user,” although Juels et al. do suggest puzzle size adjustments, as recited below;

however, Juels et al. do disclose,

- “The client puzzle protocol also allows for graceful degradation in service when an attack is mounted. The size of the puzzles can be increased as the progress of an attack advances closer to disabling the server. This enables the protocol to flex according to the scale of the attack” [column 9 lines 10-14];

Therefore, it would have been obvious for one of ordinary skill in the art at the time of the applicant’s invention to include, “wherein the step of determining an amount of accumulated work output value is performed for a specified user only during a specified time period in which accumulating work is allowed for that specified user,” in the invention as disclosed by Schuba et al. and French et al. since the client puzzle protocol is used for controlling the rate of connection buffer allocation and the likely computational capacity in order to provide graceful degradation in service when an attack is mounted (i.e. denial of service attack).

Claim 9:

Schuba et al. and French et al. disclose a method of preventing an attack on a network, as in Claim 1 above, but their combination do not disclose,

- “wherein the step of determining an amount of accumulated work output value is performed for a specified user only if the current user identity value received from the user is not found in a list of user identity values that were previously received in a specified time period,” although Juels et al. do suggest increasing puzzle size in response to an attack, as recited below;

however, Juels et al. do disclose,

- “The client puzzle protocol also allows for graceful degradation in service when an attack is mounted. The size of the puzzles can be increased as the progress of an attack advances closer to disabling the server. This enables the protocol to flex according to the scale of the attack” [column 9 lines 10-14];

Therefore, it would have been obvious for one of ordinary skill in the art at the time of the applicant's invention to include, “wherein the step of determining an amount of accumulated work output value is performed for a specified user only if the current user identity value received from the user is not found in a list of user identity values that were previously received in a specified time period,” in the invention as disclosed by Schuba et al. and French et al. since the client puzzle protocol is used for controlling the rate of connection buffer allocation and the likely computational capacity in order to provide graceful degradation in service when an attack is mounted (i.e. denial of service attack).

Claim 10:

Schuba et al. and French et al. disclose a method of preventing an attack on a network, as in Claim 1 above, but their combination do not disclose,

- “digitally signing and providing a timestamp to the user with the accumulated work output value,” although Juels et al. do suggest time stamping and usage of a secretly computed message authentication code residing as part of the other data, as recited below;

- “wherein the step of determining an amount of accumulated work output value is performed for a specified user,” although Juels et al. do suggest client puzzles, as recited below;
- “only upon: receiving the timestamp is received in a subsequent request,” although Juels et al. do suggest usage of time stamps, as recited below;
- “only upon: verifying the timestamp value,” although Juels et al. do suggest usage of time stamps, as recited below;
- “only upon: determining that the timestamp value is within an allowed range,” although Juels et al. do suggest usage of time stamps, as recited below;

however, Juels et al. do disclose,

- “This time stamp, or any other portion of seed data (SD) can be optionally authenticated with the use of a secretly computed message authentication code residing as part of the other data (OD) 530 portion of the seed data (500)” [column 19 lines 22-26];
- “The client puzzle protocol also allows for graceful degradation in service when an attack is mounted. The size of the puzzles can be increased as the progress of an attack advances closer to disabling the server. This enables the protocol to flex according to the scale of the attack” [column 9 lines 10-14];

Therefore, it would have been obvious for one of ordinary skill in the art at the time of the applicant’s invention to include, “digitally signing and providing a timestamp to the user with the accumulated work output value” and “wherein the step of determining an amount of accumulated work output value is performed for a specified user” and “only upon: receiving the timestamp is received in a subsequent request” and “only upon: verifying the timestamp value” and “only

upon: determining that the timestamp value is within an allowed range,” in the invention as disclosed by Schuba et al. and French et al. since “secretly computed message authentication code residing as part of the other data” may typically be “digitally signed and time stamped” information for verification, where a client puzzle protocol is used to control graceful degradation in service.

Claim 12:

Schuba et al. and French et al. disclose a method of preventing an attack on a network, as in Claim 1 above, but their combination do not disclose,

- “a prior user identity value and a current user identity value in a cookie provided by the user to the resource,” although Juels et al. do suggest a client puzzle protocol, as recited below;
- “wherein determining an amount of accumulated work output value to provide to the user based on a volume of data communicated between the resource and the user comprises determining the amount of accumulated work as $2^k * p$,” although Juels et al. do suggest client puzzles, as recited below;
- “where k is a number of bits of work previously performed by the user and p is a number of messages or packets communicated between the user and the resource,” although Juels et al. do suggest client puzzles, as recited below;

however, Juels et al. do disclose,

- “the “client puzzle” protocol” [column 8 line 65];

- “The client puzzle protocol also allows for graceful degradation in service when an attack is mounted. The size of the puzzles can be increased as the progress of an attack advances closer to disabling the server. This enables the protocol to flex according to the scale of the attack” [column 9 lines 10-14];

Therefore, it would have been obvious for one of ordinary skill in the art at the time of the applicant’s invention to include, “a prior user identity value and a current user identity value in a cookie provided by the user to the resource” and “wherein determining an amount of accumulated work output value to provide to the user based on a volume of data communicated between the resource and the user comprises determining the amount of accumulated work as $2^k * p$ ” and “where k is a number of bits of work previously performed by the user and p is a number of messages or packets communicated between the user and the resource,” in the invention as disclosed by Schuba et al. and French et al. since the client puzzle protocol is used for controlling the rate of connection buffer allocation and the likely computational capacity in order to provide graceful degradation in service when an attack is mounted (i.e. denial of service attack).

Claim 13:

Schuba et al. and French et al. disclose a method of preventing an attack on a network, as in Claim 1 above, but their combination do not disclose,

- “providing the accumulated work output value in a cookie sent from the resource to the user,” although Juels et al. do suggest client puzzles, as recited below;

however, Juels et al. do disclose,

- “the “client puzzle” protocol” [column 8 line 65];

Therefore, it would have been obvious for one of ordinary skill in the art at the time of the applicant's invention to include, "providing the accumulated work output value in a cookie sent from the resource to the user," in the invention as disclosed by Schuba et al. and French et al. since the client puzzle protocol is used for controlling the rate of connection buffer allocation and the likely computational capacity in order to provide graceful degradation in service when an attack is mounted (i.e. denial of service attack).

Claim 14:

Schuba et al. and French et al. disclose a method of preventing an attack on a network, as in Claim 1 above, but their combination do not disclose,

- "selectively increasing the required work threshold value for a particular user in response to congestion conditions of the resource," although Juels et al. do suggest adjusting client puzzle size in response to an attack, as recited below;

however, Juels et al. do disclose,

- "The client puzzle protocol also allows for graceful degradation in service when an attack is mounted. The size of the puzzles can be increased as the progress of an attack advances closer to disabling the server. This enables the protocol to flex according to the scale of the attack" [column 9 lines 10-14];

Therefore, it would have been obvious for one of ordinary skill in the art at the time of the applicant's invention to include, "selectively increasing the required work threshold value for a particular user in response to congestion conditions of the resource," in the invention as disclosed by Schuba et al. and French et al. since the client puzzle protocol is used for

controlling the rate of connection buffer allocation and the likely computational capacity in order to provide graceful degradation in service when an attack is mounted (i.e. denial of service attack).

Claims 21-23, 25-27, & 29-31:

Schuba et al. disclose an apparatus and a computer-readable storage medium storing one or more sequences of instructions, as in Claims 17-19 above, but they do not disclose,

- “determining the required work threshold value based on a then-current capacity of the resource,” although Juels et al. do suggest computational capacity, as recited below;
- “requiring a first user who has an accumulated work value that is greater than the required work threshold value to perform a first amount of work as a condition for accessing the resource,” although Juels et al. do suggest client puzzles, as recited below;
- “requiring a second user who has an accumulated work value that is less than or equal to the required work threshold value to perform a second amount of work as a condition for accessing the resource,” although Juels et al. do suggest client puzzles, as recited below;
- “wherein the second amount of work is greater than the first amount of work,” although Juels et al. do suggest client puzzles, as recited below;
- “determining an amount of accumulated work output value is operable for a specified user only if the current user identity value received from the user is not found in a list of user identity values that were previously received in a specified time period,” although Juels et al. do suggest client puzzles, as recited below;

- “digitally signing and providing a timestamp to the user with the accumulated work output value,” although Juels et al. do suggest time stamping and usage of a secretly computed message authentication code residing as part of the other data, as recited below;
- “determining an amount of accumulated work output value is operable for a specified user only upon: receiving the timestamp is received in a subsequent request,” although Juels et al. do suggest usage of time stamps, as recited below;
- “verifying the timestamp value,” although Juels et al. do suggest usage of time stamps, as recited below;
- “determining that the timestamp value is within an allowed range,” although Juels et al. do suggest usage of time stamps, as recited below;

however, Juels et al. do disclose,

- “the rate of connection buffer allocation and the likely computational capacity of one or more attacking clients 110 can be used to select the computational size of a particular tasks when operating in a defensive mode” [column 7 lines 29-33];
- “The client puzzle protocol also allows for graceful degradation in service when an attack is mounted. The size of the puzzles can be increased as the progress of an attack advances closer to disabling the server. This enables the protocol to flex according to the scale of the attack” [column 9 lines 10-14];
- “inside each correct sub-puzzle solution, and comparing the time stamp (DT) with the current time to check that the (sub)puzzle has not yet expired” [column 19 lines 20-22];

- “This time stamp, or any other portion of seed data (SD) can be optionally authenticated with the use of a secretly computed message authentication code residing as part of the other data (OD) 530 portion of the seed data (500)” [column 19 lines 22-26];

Therefore, it would have been obvious for one of ordinary skill in the art at the time of the applicant's invention to include, “determining the required work threshold value based on a then-current capacity of the resource” and “requiring a first user who has an accumulated work value that is greater than the required work threshold value to perform a first amount of work as a condition for accessing the resource” and “requiring a second user who has an accumulated work value that is less than or equal to the required work threshold value to perform a second amount of work as a condition for accessing the resource” and “wherein the second amount of work is greater than the first amount of work” and “determining an amount of accumulated work output value is operable for a specified user only if the current user identity value received from the user is not found in a list of user identity values that were previously received in a specified time period” and “digitally signing and providing a timestamp to the user with the accumulated work output value” and “determining an amount of accumulated work output value is operable for a specified user only upon: receiving the timestamp is received in a subsequent request” and “verifying the timestamp value” and “determining that the timestamp value is within an allowed range,” in the invention as disclosed by Schuba et al., since the client puzzle protocol is used for controlling the rate of connection buffer allocation and the likely computational capacity in order to provide graceful degradation in service when an attack is mounted (i.e. denial of service attack).

Response to Arguments

9. Applicant's arguments with respect to claims 1-15 have been considered but are moot in view of the new ground(s) of rejection as necessitated by the applicant's amendments.
10. Applicant's arguments, see pages 16-18, filed 05/23/2008, with respect to the rejection(s) of claim(s) 16 under 35 U.S.C. 103(a) have been fully considered and are persuasive. Therefore, the rejection has been withdrawn. However, upon further consideration, a new ground(s) of rejection is made in view of newly found prior art reference(s).
11. Applicant's arguments filed 05/23/2008 with respect to Claims 17-31 have been fully considered but they are not persuasive.
- The applicant's remarks with respect to Claims 17-31 have been carefully considered, however, the currently written claims language for these limitations are broad enough to be interpreted as reading on the prior art of record as indicated in the rejections above;
 - o The examiner notes that one interpretation of the "accumulated work value" would be the overall or current work value that the client/user is currently generating through its calculations of the "puzzle" solution; that is, as is currently claimed, the applicant's limitations appear to be directed towards the same "client puzzle" scheme as the prior art; further claim limitations, such as the amendments that were made to independent Claim 1 above, are not found in Claims 17-31

Conclusion

12. Any inquiry concerning this communication or earlier communications from the examiner should be directed to Examiner Oscar Louie whose telephone number is 571-270-1684. The examiner can normally be reached Monday through Thursday from 7:30 AM to 4:00 PM.

If attempts to reach the examiner by telephone are unsuccessful, the examiner's supervisor, Nasser Moazzami, can be reached at 571-272-4195. The fax phone number for Formal or Official faxes to Technology Center 2100 is 571-273-8300.

Information regarding the status of an application may be obtained from the Patent Application Information Retrieval (PAIR) system. Status information for published applications may be obtained from either Private PAIR or Public PAIR. Status information for unpublished applications is available through Private PAIR only. For more information about the PAIR system, see <http://pair-direct.uspto.gov>. Should you have questions on access to the Private PAIR system, contact the Electronic Business Center (EBC) at 866-217-9197 (toll-free). If you would like assistance from a USPTO Customer Service Representative or access to the automated information system, call 800-786-9199 (IN USA OR CANADA) or 571-272-1000.

OAL
07/25/2008

/Nasser G Moazzami/
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